# Sparse Merkle Trees

Introducing the concept and benchmarking libraries available in Rust

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### What is the problem

In decentralization where we rely on untrusted parties, exchanges requires **integrity** and **authenticity**, at a very **big scale**.

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**Authenticity**: The property that data originated from its purported source.

Solved using **cryptographic signatures** (RSA, ECDSA), generated using a *secret key* only the owner has, can be verified by anyone using the *public key* associated to the secret key.

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#### **Situation**

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Hacker Pizza's WIFI alters the data she receives

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#### **Situation**

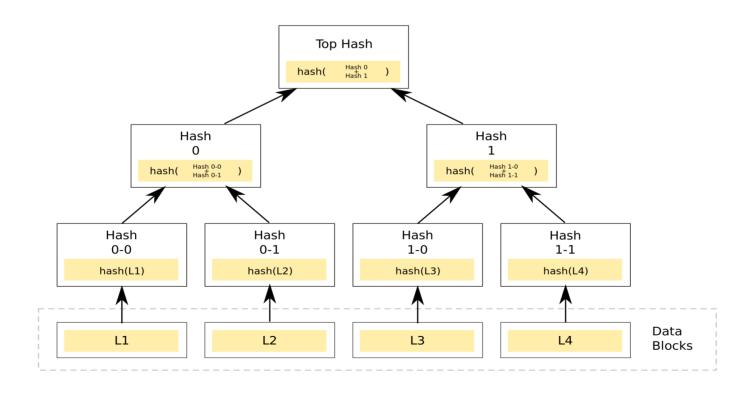
Alice paid a pizza 10 BTC in "Hacker Pizza" using her smartphone, want to check if transaction accepted.

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#### Only solution to protect herself

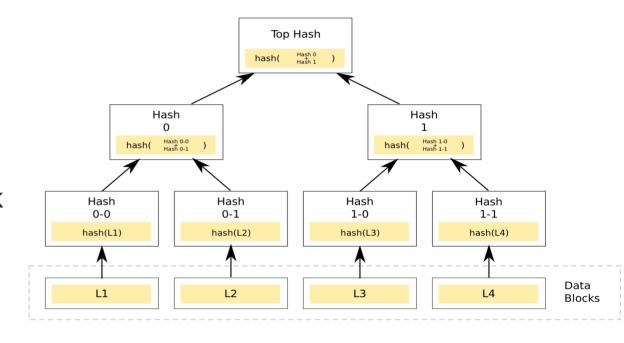
Hash all the transactions, compare with hash given in block header, then verify the block header's signature is correct.

Hash large data, pieces by pieces, without compromises on integrity.



Alice wants to verify the transaction L2

- Hash the transaction: H0-1
- Ask H1 and H0 0, verify signatures
- H0 = Hash(H0-0, H0-1)
- TopHash = Hash(H0, H1)
- Compare Top hash to the one in block header
- Verify signature of the hash in block header



Merkle Proof: All the data needed to verify a leaf of the tree

For Alice's tx: Hash(L2) + H0-0 + H1 (concatenated)

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Have hash of a modified datastore with N elements requires O(log(N)) hashes operations.

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```
Add new data to the ledger: F' = F XOR Hash(new_data)
Remove data off the ledger: F' = F XOR Hash(rm_data)
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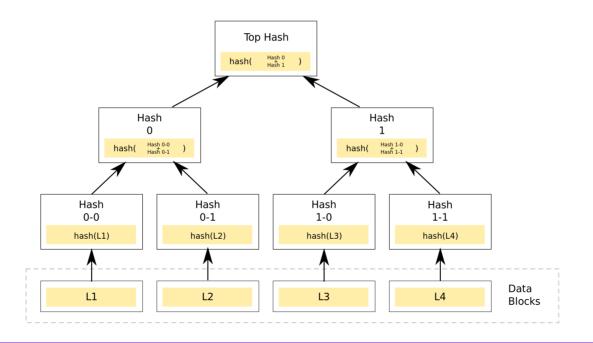
```
Add new data to the ledger: F' = F XOR Hash(new_data)
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```

Simple, fast, incremental

Not a hash function, not suited for integrity checks If A XOR B = 0, then C XOR (A XOR B) = C

# Why we want it in Massa So let's use Merkle Trees! Why Sparse?

What if we want to add a new data between L1 and L2?



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Proof of **non-inclusion** 

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Integrity check over the whole data

No compromises

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Note: Incremental hash was also a different posibility, look at the discussion on Github to find more details on why SMT was chosen.

Or ask to Varun, he explains it very well:-)

```
struct SparseMerkleTree<H: Hasher, D: Database> { ... }
```

```
struct SparseMerkleTree<H: Hasher, D: Database> { ... }

trait Database {
   fn get(...)
   fn put(...)
   fn remove(...)
}
```

```
struct SparseMerkleTree<H: Hasher, D: Database> { ... }
trait Database {
 fn get(...)
 fn put(...)
 fn remove(...)
trait Hasher {
 fn new(...)
  fn update(...)
 fn finalize(...)
```

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Very implementation-dependant

#### **Benchmarks**

Implemented Blake3Hasher, MemoryStorage, RockSdbStorage

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Framework	Last updated	Stars on Github
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#### **Benchmarks**

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cw-merkle-tree was ignored as it's too tied to CosmWasm smart contract framework

#### **Benchmarks**

On MemoryStore (storage in RAM)

```
monotree/memstore+blake3
  time: [17.473 μs 17.619 μs 17.770 μs]

sparse-merkle-tree/memstore+blake3
  time: [119.37 μs 120.63 μs 122.11 μs]

lsmtree/memstore+blake3
  time: [25.587 μs 25.768 μs 25.952 μs]
```

#### **Benchmarks**

On RocksDB (storage on the disk)

```
monotree/rocksdb+blake3

time: [153.27 μs 155.79 μs 158.37 μs]

sparse-merkle-tree/rocksdb+blake3

time: [1.3083 ms 1.3135 ms 1.3190 ms]

lsmtree/rocksdb+blake3

time: [248.28 μs 249.85 μs 251.47 μs]
```

# Read / Write operations benchmark

On MemoryStore (storage in RAM)

```
monotree/memstore+blake3/read
   time: [2.7194 μs 2.7374 μs 2.7564 μs]

lsmtree/memstore+blake3/read
   time: [174.66 ns 178.09 ns 181.79 ns]

monotree/memstore+blake3/write
   time: [14.213 μs 14.318 μs 14.423 μs]

lsmtree/memstore+blake3/write
   time: [24.849 μs 25.431 μs 26.049 μs]
```

# Read / Write operations benchmark

On RocksDB (storage on the disk)

```
monotree/rocksdb+blake3/read
   time: [10.370 μs 10.646 μs 10.938 μs]

lsmtree/rocksdb+blake3/read
   time: [581.90 ns 609.57 ns 638.72 ns]

monotree/rocksdb+blake3/write
   time: [150.39 μs 161.96 μs 172.92 μs]

lsmtree/rocksdb+blake3/write
   time: [233.95 μs 239.78 μs 245.64 μs]
```

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Fully featured already

Simple but efficient

Can be maintained by our own means

# Some links

Benchmark code

Article on a performance-oriented SMT implementation

How Merkle trees is used in Bitcoin

Github discussions about implementing SMT in Massa

Why use binary trees over trees with more children

Libra whitepaper, contains optimizations for SMT